

Securing Inter-process Communications in SELinux

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Introduction

- Inter-process communication (IPC)
 - method of information exchange between processes
 - local machine, or network
 - used to share data, state, and synchronization
- Numerous forms for various use cases
 - files, message queues, shared memory, etc.
- Each has documented benefits
 - throughput, latency, ease of use, portability
- Linux, SELinux supports BSD, POSIX, and SysV IPC
 - don't worry, no Mach Messages

The Problem

- Little to no IPC security documentation
 - current docs (man pages, books) don't address information-flow security
- How do I know if a denial is my problem or yours?
 - control options, write implies read, etc.
- Which type of IPC should I use?
 - RTFM
 - rely on intuition and experience
- Isn't there an easier way to answer these questions?

Objectives

- Security Objectives
 - provide access control
 - ensure proper directionality
 - facilitate revocation of access
 - protect integrity of both data and processes
- System objectives must also be considered
 - functional requirements
 - performance, usage
 - security goals
 - risks and threat model

Security Objectives (cont)

- Fine-grained access controls
 - fine-grained mediation of access on most IPC
 - controls read, write, control options, etc
 - overloaded, userspace supported the exception
- Revocation
 - must be able to rescind access previously granted
 - SELinux supports access time revocation on most IPC
 - “memory based” IPC is the exception
 - attributed to kernel architecture, unlikely to change

Security Objectives (cont)

- Unidirectional information flow
 - not just for guard/CDS use
 - pipes, FIFO, datagrams
 - hard to achieve
 - guarantees of delivery almost always necessary
 - numerous backchannels exist, such as blocking
- Integrity
 - SELinux provides integrity
 - fine-grained access control

Pipes & Unix Domain Sockets

- Pipes & FIFOs (named/unnamed Pipes)
 - SELinux *can* enforce directionality but...
 - non-blocking storage channel via partial writes
 - FIFOs also have flow through the file system
 - access time revocation supported
- Unix Domain Sockets
 - creation bitmap
 - channels through files
 - access time revocation supported

Shared Memory & Message Queues

- Shared memory
 - access controls fine-grained
 - possible to enforce directionality
 - synchronization in the unidirectional case?
 - permissions migrated
- Message queues
 - access controls eliminate the large back channels
 - blocking back channel exists
 - helper process eliminates “bitmap” attack vector

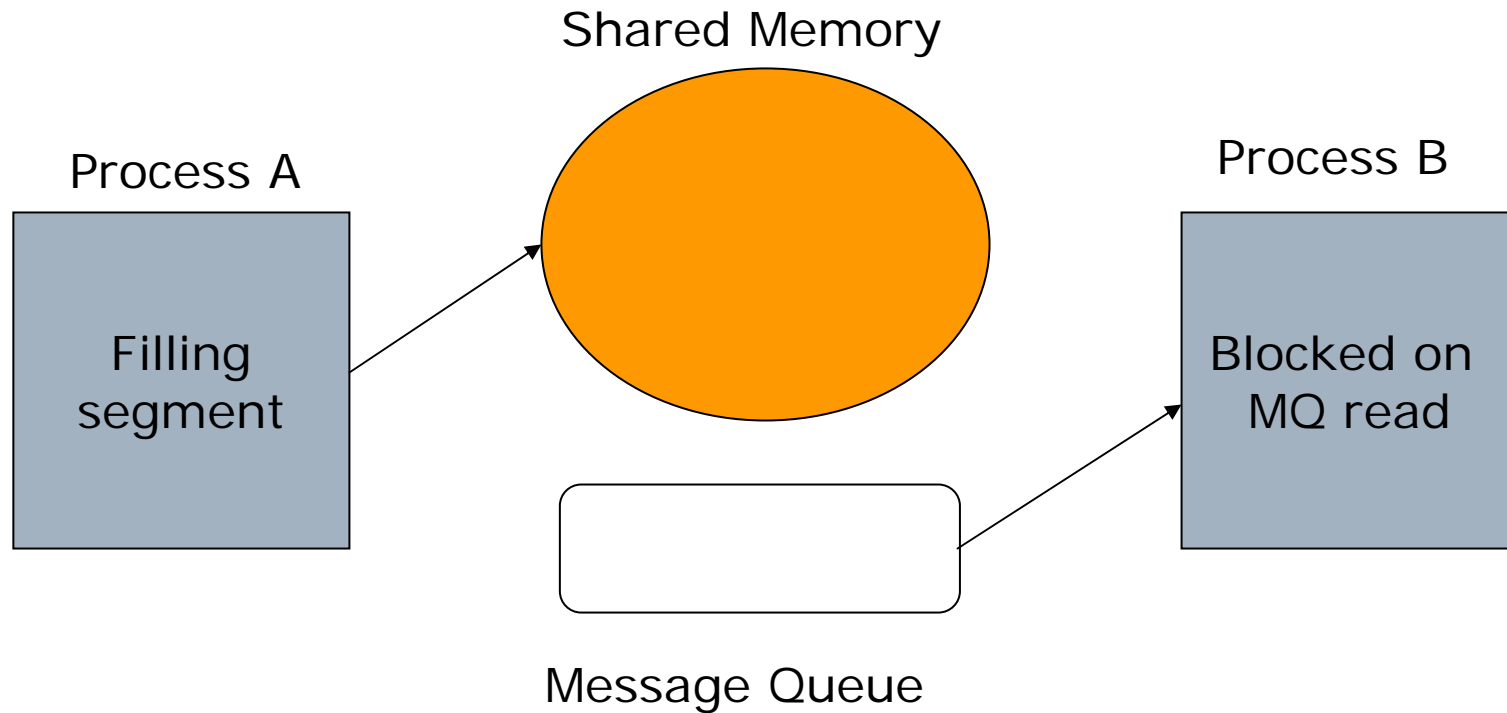
Building a Secure IPC Mechanism

- Unidirectional shared memory needs synchronization or control channel
- Message queues suit that need
- Shared memory acts as large bandwidth data channel
 - **forward only**
- Message queues act as control channel
 - **small, noisy back channel via blocking**

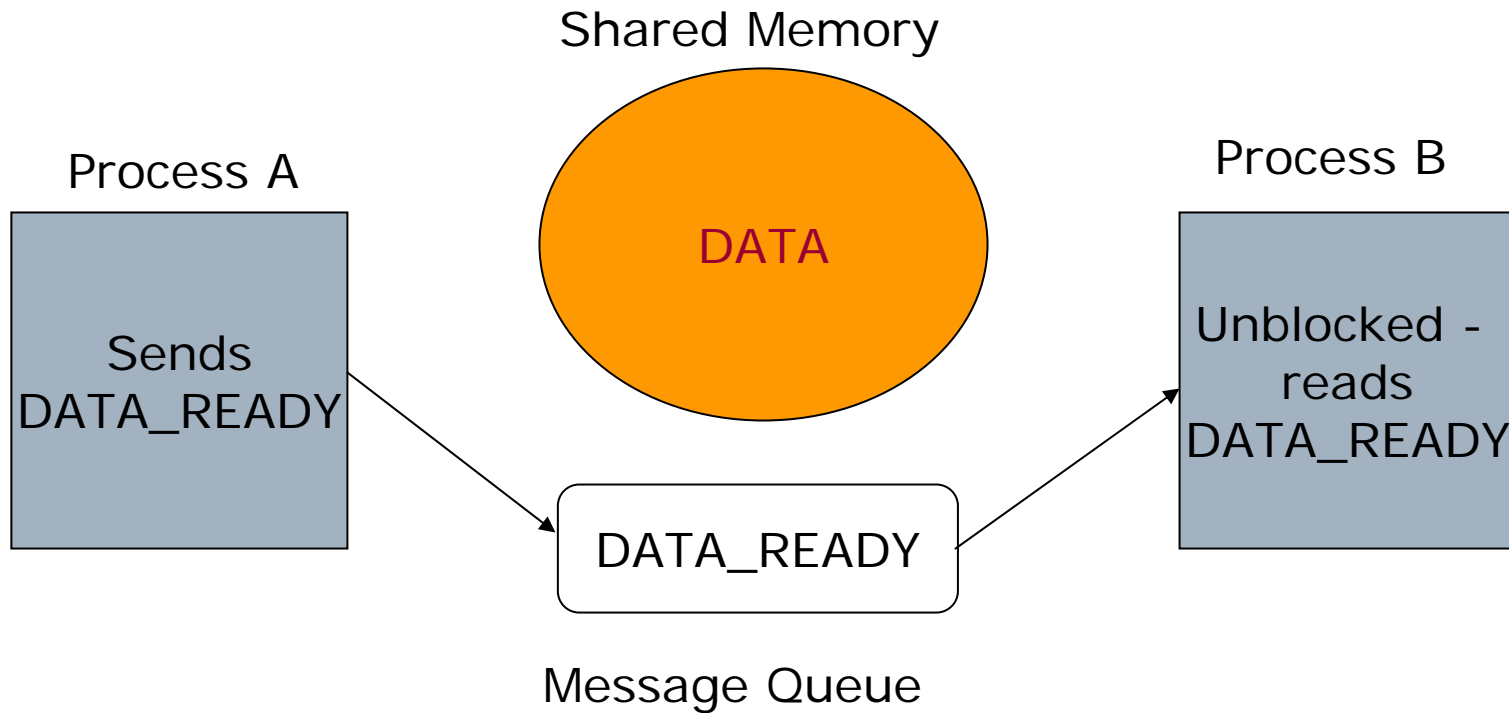
Using IPC Securely

- Helper process
 - starts, creates queue, terminates
 - eliminates bitmapping attack vector
 - neither process can create queues
 - queue retains the type of this process
- Participating processes
 - associate with the queue created above
 - process A creates a shared memory segment and attaches
 - process B attaches

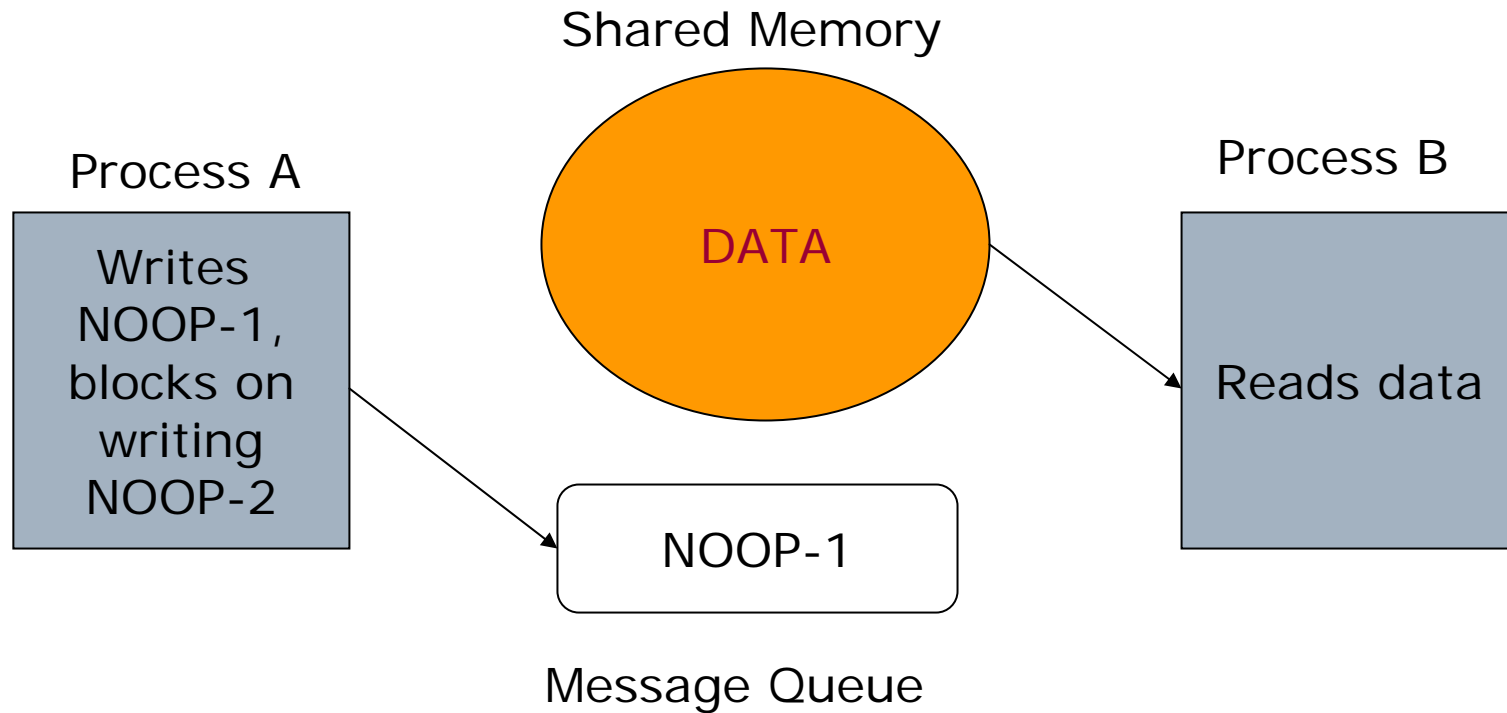
Using IPC Securely- not drawn to scale



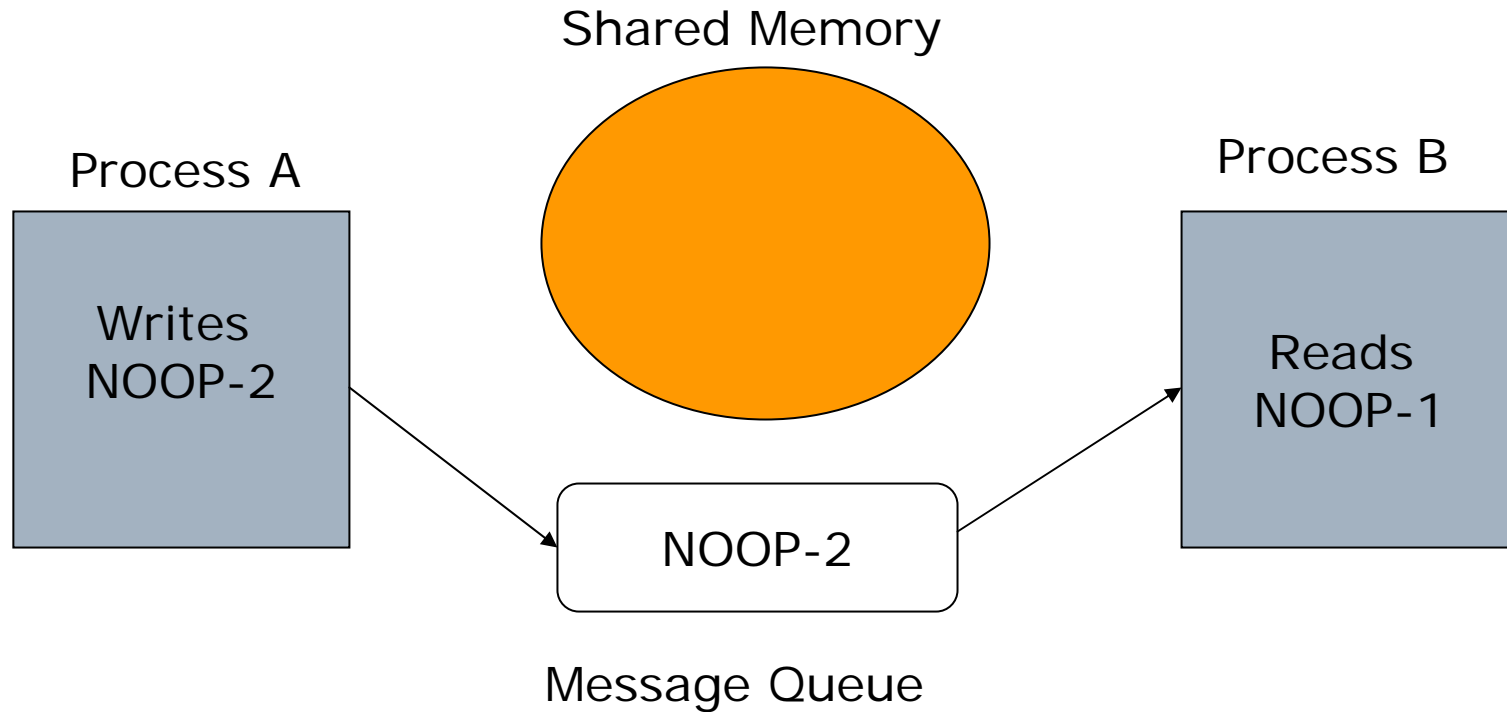
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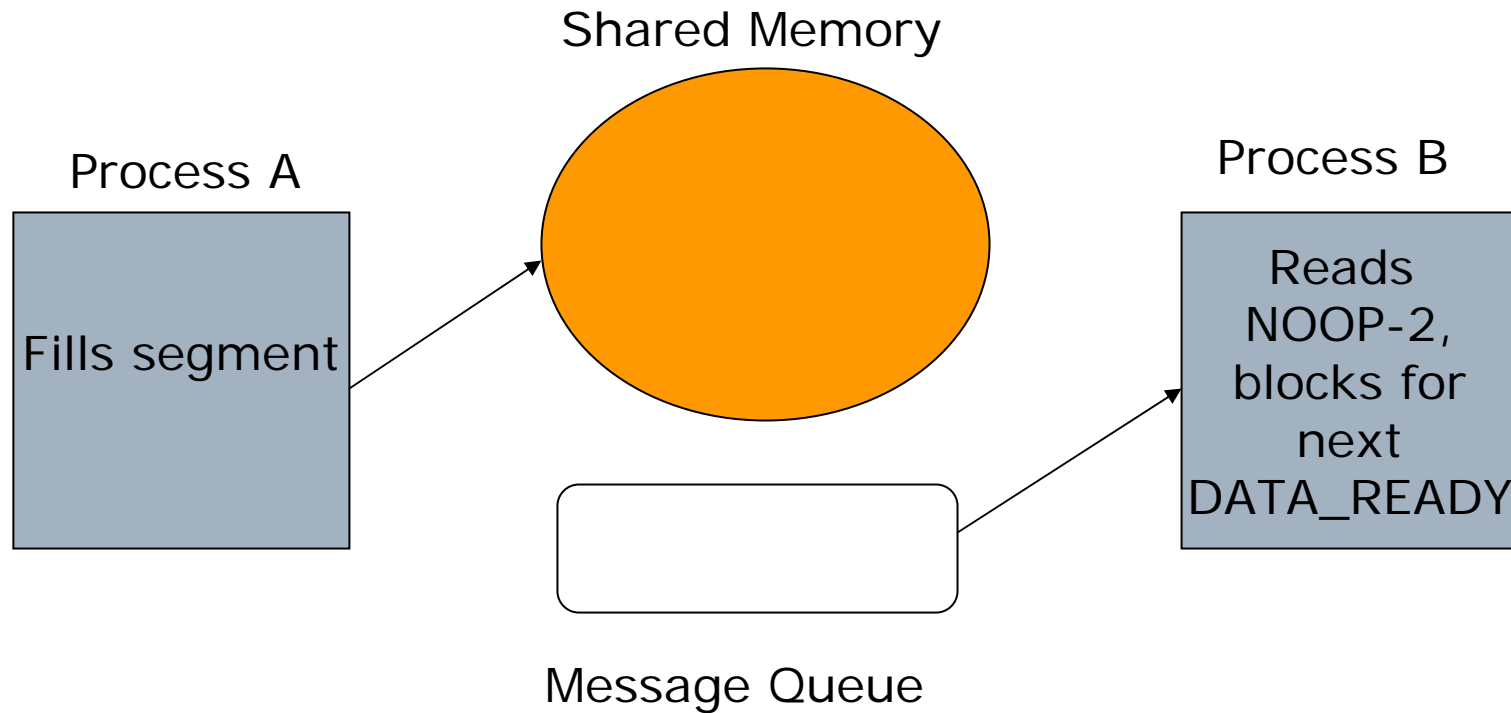
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Using IPC Securely- not drawn to scale



Using IPC Securely

1. Wash
2. Rinse
3. Repeat

Conclusions

- IPC is not unidirectional by definition
 - SysV, POSIX, BSD APIs compliance requires backchannels
- In SELinux, message queues and shared memory
 - a large bandwidth forward channel
 - a limited bandwidth back channel
 - enforced with a MAC policy
- C library, Java and Mono bindings, and policy available
 - <http://oss.tresys.com/projects/sipc>

Questions?

<http://oss.tresys.com/projects/sipc>